

# Almost Linear Constant-Factor Sketching for $\ell_{1}$ and Logistic Regression 

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## Setting: Logistic regression

- Input: Dataset $\left\{x_{1}, \ldots, x_{n}\right\}, x_{i} \in \mathbb{R}^{d}$ with labels $y_{i} \in\{-1,1\}$.

■ We search for the empirical risk minimizer $\beta \in \mathbb{R}^{d}$ minimizing

$$
f(\beta)=\sum_{i=1}^{n} \ln \left(1+\exp \left(-y_{i} x_{i} \beta\right)\right) .
$$

## Logistic regression for massive data

## Problems:

- Too much data to store in the working memory
- Limited access to data: data streams
- Data is given in pieces or updated dynamically (turnstile streams, vertically distributed)


## Massive data analysis

## Sketch and solve paradigm



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Sketch and solve paradigm


Canonical approach
11 Data reduction $X \rightarrow \Pi X$ (fast linear sketch), where $|\Pi X| \ll|X|$
2 Time- and space efficient calculations on $\Pi X$
3 Approximation guarantee: solution is close to optimal

## Sketching matrix

$$
\Pi=\left(\begin{array}{c}
S_{0} \\
S_{1} \\
\vdots \\
S_{h_{\max }}
\end{array}\right)
$$

- Idea: Subsample the points at different rates: for each entry randomly choose its level $h \in \mathbb{N}$ at rate $2^{-h}$
$-S_{h_{\max }}$ is a uniform sample and $S_{0}$ is a CountMin-Sketch of the full data;


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- $S_{h_{\max }}$ is a uniform sample and $S_{0}$ is a CountMin-Sketch of the full data;
- To reflect the rarity of entries, buckets of high levels (with few elements) get a high weight;
■ Sketch the points at each level: inside the levels we use the CountMin-Sketch, i.e., we hash points into buckets uniformly at random and add up entries in the same bucket.


## Analysis

Fix $\beta$ and set $z=X \beta$, prove that with high probability following holds:

- Contraction bounds: $f(\Pi z) \geq(1-\varepsilon) f(z)$ (apply a net argument to show that this holds for any $z^{\prime}=X \beta^{\prime}$ );
- Dilation bounds: $f(\Pi z) \leq k f(z)$.


## Analysis

Idea: Divide the entries of $z$ into weight classes
$W_{q}^{+}=\left\{i \left\lvert\, 2^{-q} \geq \frac{z_{i}}{\|z\|_{1}}>2^{-(q+1)}\right.\right\}$.

- For any 'relevant' weight class $W_{q}^{+}$there is a level where the contribution of $W_{q}^{+}$is preserved
- The expected contribution of any weight class $W_{q}^{+}$to any level is at most $\left\|W_{q}^{+}\right\|_{1}$ and $W_{q}^{+}$contributes to at most $k$ levels


$$
\begin{equation*}
h=h_{m} \tag{m}
\end{equation*}
$$

## Main result

## Theorem 1

There is a distribution over sketching matrices $\Pi \in \mathbb{R}^{r \times n}$ such that with constant probability
$1 \Pi X$ has $r=O\left(\mu d^{1+c} \ln (n)^{2+4 c}\right)$ rows, can be computed in time $O(d \ln (n) \mu \cdot \mathrm{nnz}(X))$ and yields an $O(1)$ approximation.
2 $\Pi X$ has $r=\mu^{2}\left(\varepsilon^{-1} \ln (n) d\right)^{O\left(\varepsilon^{-1}\right)}$ rows, can be computed in time $O(\mathrm{nnz}(X))$ and yields an $1+\varepsilon$ approximation.
$\mu=\sup _{\beta \in \mathbb{R}^{d} \backslash\{0\}} \frac{\sum_{x_{i} \beta>0}\left|x_{i} \beta\right|}{\sum_{x_{i} \beta<0}\left|x_{i} \beta\right|}$
$\mathrm{nnz}(X)=$ number of non-zero entries of $X$.

## Other target functions

We get similar results for $\ell_{1}$-regression

$$
f(X \beta)=\|X \beta-y\|_{1}
$$

and logistic regression with variance-based regularization

$$
\begin{aligned}
f(X \beta) & =\frac{1}{n} \sum_{i=1}^{n} \ell\left(x_{i} \beta\right)+\frac{\lambda}{2 n} \sum_{i=1}^{n} \ell\left(x_{i} \beta\right)^{2}-\frac{\lambda}{2}\left(\frac{1}{n} \sum_{i=1}^{n} \ell\left(x_{i} \beta\right)\right)^{2} \\
& =\mathbb{E}(\ell(x \beta))+\frac{\lambda}{2} \cdot \operatorname{Var}(\ell(x \beta)) .
\end{aligned}
$$

